## Big O Notation

## Last updated 1/30/23

These slides introduce big-O notation

- Big O Notation
- Describes the limit of a function
- Provides an asymptotic upper (lower) bound of the function
- In programming
- Mathematical tool to measure the cost of an algorithm
- Cost can be
- Operations to execute (not time)
- Memory Needed
- Energy required to complete
- Rules
- Create a function to represent the value (cost) you want to measure
- Remove all terms except the primary term
- Ignore constants of proportionality (multiplying constants)
- Determine the limit of the function as the input reaches a specific value (usually infinity)

$$
\begin{array}{ll}
\cos =4 n^{2}+2 n+4095 & \\
\operatorname{cost} \approx 4 n^{2} & \text { primary term }(n \rightarrow \infty) \\
\cos t \approx n^{2} & \text { no constants } \\
\operatorname{cost}=O\left(n^{2}\right) &
\end{array}
$$

- Examples
- Operations to read an individual array value $\rightarrow O(1)$
- Operations to print an entire 1-d array $\rightarrow \mathrm{O}(\mathrm{n})$
- Operations to print a 2-d array $\rightarrow$ O(n)
- Note: $n$ here is defined as the number of elements
- Operations to print all pairs of values in a 1-d array $\rightarrow$ $O\left(n^{2}\right)$
- Operations to calculate Fibonacci sequence recursively $\rightarrow$ $O\left(2^{n}\right)$
- Operations to calculate Fibonacci sequence with a for loop $\rightarrow \mathrm{O}(\mathrm{n})$


## Big O Notation

## - Relative Complexity (growth)

## Big-O Complexity Chart



Elements
src: bigocheatsheet.com

## Big O Notation

## - Common Structures

Common Data Structure Operations

| Data Structure | Time Complexity |  |  |  |  |  |  |  | Space Complexity <br> Worst |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Average |  |  |  | Worst |  |  |  |  |
|  | Access | Search | Insertion | Deletion | Access | Search | Insertion | Deletion |  |
| Array | $\theta(1)$ | $\theta(\mathrm{n})$ | $\theta(\mathrm{n})$ | $\theta(\mathrm{n})$ | O(1) | $O(n)$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ |
| Stack | $\theta(\mathrm{n})$ | $\theta(\mathrm{n})$ | $\theta(1)$ | $\theta(1)$ | $0(\mathrm{n})$ | $0(n)$ | 0 (1) | 0 (1) | $0(\mathrm{n})$ |
| Queue | $\theta(\mathrm{n})$ | $\theta(\mathrm{n})$ | $\theta(1)$ | $\theta(1)$ | $O(n)$ | $O(n)$ | O(1) | O(1) | $0(n)$ |
| Singly-Linked List | $\theta(\mathrm{n})$ | $\theta(\mathrm{n})$ | $\theta(1)$ | $\theta(1)$ | $0(n)$ | $O(n)$ | O(1) | O(1) | $0(n)$ |
| Doubly-Linked List | $\theta(\mathrm{n})$ | $\theta(\mathrm{n})$ | $\theta(1)$ | $\theta(1)$ | $0(\mathrm{n})$ | $O(n)$ | 0(1) | O(1) | $0(\mathrm{n})$ |
| Skip List | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(n \log (n))$ |
| Hash Table | N/A | $\theta(1)$ | $\theta(1)$ | $\theta(1)$ | N/A | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $O(n)$ |
| Binary Search Tree | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $O(n)$ | $0(n)$ | $0(n)$ | $0(\mathrm{n})$ | $0(n)$ |
| Cartesian Tree | N/A | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | N/A | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ |
| B-Tree | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (\mathrm{n}) \mathrm{)}$ | $0(\mathrm{n})$ |
| Red-Black Tree | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (n))$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (n))$ | $O(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (\mathrm{n}))$ | $0(\log (n))$ | $0(\mathrm{n})$ |
| Splay Tree | N/A | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | N/A | $0(\log (\mathrm{n}) \mathrm{)}$ | $O(\log (\mathrm{n})$ ) | $O(\log (n))$ | $0(\mathrm{n})$ |
| AVL Tree | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (n))$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $O(\log (n))$ | $0(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (\mathrm{n}) \mathrm{)}$ | $0(\log (n))$ | $0(n)$ |
| KD Tree | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $\theta(\log (\mathrm{n}) \mathrm{)}$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ | $0(\mathrm{n})$ |
|  |  |  |  |  |  |  | : bigochea | atsheet.com |  |

Note: If n is not in the function, we get $\mathrm{O}(1)$

